Monday: Linear Diophantine equations and cancellation laws

\mathcal{L} are discrete expansions equations

A *Diophantine equation* is an algebraic equation (e.g. $ax^2 + bx + cxy = d$) in which the coefficients (a, b, c and d) are integers, and for which we seek integer solutions x and y. We will consider the special case c and d) are integers, and for which we see the second considerations α and β . We will consider the special case of linear Diophantine equations, which are of the form

$$
ax + by = c,\tag{*}
$$

where a, b, c $\in \mathbb{Z}$: we seek all integers x and y satisfying the equation (*). Of course, if x and y were allowed to be real numbers, then (∗) would be the equation of a straight line: we ask when this straight line intersects the lattice of points $\mathbb{Z}^2 = \{ (x, y) : x, y \in \mathbb{Z} \}$. In general, a straight line could intersect \mathbb{Z}^2 in no points (e.g. $y = x + \sqrt{2}$), in one point (e.g. $y = \sqrt{2}x$, which intersects \mathbb{Z}^2 only at the point $(0, 0)$) or infinitely often (e.g. $y = x$). When we insist on integer coefficients only the first and the third possibilities occur.

we will assume that $a, b \neq 0$. Put $d = \gcd(a, b)$. We know that $d | a$ and $d | b$, so for any $x, y \in \mathbb{Z}$ we have $d | a x + bu$. Thus if $(*)$ has a solution we must have $d | c$; if $d k c$ then no solution is possible. $d \mid ax + by$. Thus if (*) has a solution, we must have $d \mid c$: if $d \nmid c$ then no solution is possible.

 \mathbf{b} and side would be even but the right hand side would be odd hand side would be even but the right hand side would be odd.

So suppose that d | c, in other words $c = dq$ for some q. Now, we know that there exist $x_d, y_d \in \mathbb{Z}$ with $d = ax_d + by_d$. Multiplying by q we get $dq = ax_dq + by_dq$, i.e. $c = a(x_dq) + b(y_dq)$. Thus (x_dq, y_dq) is a solution of $(*).$

Example 2. Find a solution to the equation $4x + 7y = 13$.

What about the general solution? What happens if we try to prove the solution is unique?

Suppose that (x, y) and (x', y') are solutions. Then we have

$$
ax + by = c = ax' + by',
$$

so $a(x - x') + b(y - y') = 0$, or $a(x - x') = b(y' - y)$. Does this imply that $x - x' = y' - y = 0$? No, it
only implies that the number $a(x - x')$ is a common multiple of a and b. If m is any common multiple of only implies that the number $a(x - x')$ is a common multiple of a and b. If m is any common multiple of a and b, say $m = ra = sb$, then we can put $x' = x - r$, $y' = y + s$ to get

$$
a(x - x') + b(y - y') = a(x - (x - r)) + b(y - (y + s)) = ar - bs = m - m = 0,
$$

as required. So the general solution is given by $x = qx_d - m/a$, $y = qy_d + m/b$, where m is a common multiple of a and b. Note that m is a common multiple of a and b if and only if $lcm(a, b) \mid m$. So the general solution is $x = qx_d - tl/a$, $y = qy_d + tl/b$, where $l = \text{lcm}(a, b)$ and $t \in \mathbb{Z}$. Also, from Assignment 7, Question 1 we know that $ld = ab$, so $l/a = b/d$ and $l/b = a/d$. Combining these facts we have the following theorem.

Theorem 3. Let $a, b, c \in \mathbb{Z}$ with $a, b \neq 0$. Put $d = \gcd(a, b)$, and fix $x_d, y_d \in \mathbb{Z}$ with $d = ax_d + by_d$. Then the equation $ax + by = c$ has no integer solutions if $d \nmid c$, and has the general solution $x = \frac{c}{d}x_d - t_d^b$,
 $y = \frac{c}{d}y_d + t_d^a$ for $t \in \mathbb{Z}$ if $d \mid c$. $y = \frac{c}{d}y_d + t\frac{a}{d}$ for $t \in \mathbb{Z}$ if $d \mid c$.

Example 4. Find the general solution of the Diophantine equation $4x + 7y = 13$.

Example 5. Find the general solution of the Diophantine equation $6x - 15y = 27$.

Cancellation laws

In Z we have two cancellation laws: "if $a + c = b + c$ then $a = b$ " and "if $ac = bc$ and $c \neq 0$ then $a = b$ ". The first is easy to prove from the axioms: if $a + c = b + c$ then we have

$$
(a + c) + (-c) = (b + c) + (-c)
$$

\n
$$
a + (c + (-c)) = b + (c + (-c))
$$

\n
$$
a + 0 = b + 0
$$

\n
$$
a = b
$$
\n
$$
(definition of -c)
$$

\n
$$
(definition of 0)
$$

 \mathbb{Z} and into Q, and multiply both sides by $\frac{1}{c}$, but that relies on other things, not on the axioms for the integers. To get the cancellation law from the axioms alone, we would have to do a little work. One way integers. To get the cancellation law from the axioms alone, we would have to do a little work. One way to prove it would be to prove by induction that the result holds for all $c \in \mathbb{N}$, and then extend the result to negative values of c. We will leave this as an exercise to negative values of c. We will leave this as an exercise.

Tuesday: Review

Thursday: Class Test

Friday: Congruence Modulo ⁿ

When we considered equivalence relations we had as an example the relation \sim on Z defined by declaring that for $m, n \in \mathbb{Z}$ we have

$$
m \sim n \iff 5 \mid m - n.
$$

We showed that ∼ is an equivalence relation. This relation is called *congruence modulo 5*. In general, if $n \in \mathbb{N}$ we say that a and b are congruent modulo n if $n \mid a - b$: we write this relation $a \equiv b \pmod{n}$. This relation is an equivalence relation for every $n \in \mathbb{N}$. The set of equivalence classes is called the *integers* modulo n, written \mathbb{Z}_n . For $a \in \mathbb{Z}$, we call the equivalence class of a under congruence modulo n the *congruence class* of a, and denote it by \overline{a} .

Example 6. Fix $n = 5$. Find $\overline{0}$, $\overline{1}$, $\overline{10}$ and $\overline{16}$.

Lemma 7. Let $a, b \in \mathbb{Z}$, $n \in \mathbb{N}$. Then $a \equiv b \pmod{n}$ iff a and b give the same remainder when divided by n.

From this we know that there are exactly n congruence classes in \mathbb{Z}_n , because there are n possible remainders $0, 1, \ldots, n-1$. So we have

$$
\mathbb{Z}_n = \{\overline{0}, \overline{1}, \overline{2}, \ldots, \overline{n-1}\}.
$$

The set \mathbb{Z}_n inherits some properties from \mathbb{Z} . The most important is that we can define addition and multiplication on \mathbb{Z}_n in a natural way.

Definition. We define the operations $+_{n}$ and \cdot_{n} on \mathbb{Z}_{n} by declaring that, for $a, b \in \mathbb{Z}$,

$$
\overline{a} +_{n} \overline{b} = \overline{a+b}
$$
 and $\overline{a} \cdot_{n} \overline{b} = \overline{ab}$.

Of course we can write down any definition we like: we could define *n* to be the least positive solution of the equation $x = x + 1$ For this definition to make sense we have to make sure that the operations are well-defined. For example, with $n = 5$ consider finding $\overline{3} + \overline{7}$ and finding $\overline{18} + \overline{22}$. We have well-defined. For example, with $n = 5$, consider finding $3 + 5$ *l* and finding $18 + 522$. We have

$$
\overline{3} +_5 \overline{7} = \overline{3 + 7} = \overline{10} = \overline{0}
$$
 and $\overline{18} +_5 \overline{22} = \overline{18 + 22} = \overline{40} = \overline{0}$.

Thus we get the same answer both times. This is just as well, because $3 = 18$ and $7 = 22$, so we were doing the same sum in both cases.

For the definitions of \overline{a} and \overline{b} and \overline{a} and \overline{b} and \overline{b} and \overline{b} is \overline{b} then we get the same answer when we work out $\overline{a} +_n b$ and when we work out $a' +_n b'$, and similarly for \cdot_n . In other words, we must show that if $a \equiv a' \pmod{n}$ and $b \equiv b' \pmod{n}$ then $a + b \equiv a' + b' \pmod{n}$ and $ab \equiv a'b'$
(mod n) \mathcal{N} and \mathcal{N} .

Lemma 8. Let $a, b, a', b' \in \mathbb{Z}$, $n \in \mathbb{N}$. If $a \equiv a' \pmod{n}$ and $b \equiv b' \pmod{n}$ then $a + b \equiv a' + b' \pmod{n}$
and $ab \equiv a'b' \pmod{n}$ and $ab \equiv a'b' \pmod{n}$.

 T o understand what we have done we should see an example where the operations would not be well

Example 9. Partition $\mathbb Z$ into the three sets $\Omega = \{A, B, C\}$

$$
A = \mathbb{N}
$$

\n
$$
B = \{0\}
$$

\n
$$
C = \{-n : n \in \mathbb{N}\}.
$$

We try to define addition $+^{\prime}$ and multiplication \cdot^{\prime} by taking a representative from the two classes we are adding a representative from the two classes we are adding a representative from the consulence class of For example we have $A \cdot B = B$ because $n \cdot 0 = 0 \in B$ for every $n \in A$, and $A \cdot C = C$ because $m \cdot (-n) = -(mn) \in C$ for every $m \in A - n \in C$. However, addition is not well-defined; when we true $m \cdot (-n) = -(mn) \in C$ for every $m \in A$, $-n \in C$. However, addition is **not** well-defined: when we try to find $A +^{\prime}C$ we could get the answer A (for example by choosing the representatives 8 and -3), B (e.g. by choosing 6 and -6) or C (e.g. by choosing 5 and -12). The answer we get depends not just on the classes but on which representative of the classes we choose.

What can we say about arithmetic modulo n? We know that the operations $+_n$ and \cdot_n are commutative and associtive, and \cdot_n distributes over \cdot_n . To show the last one, let $a, b, c \in \mathbb{Z}$. Then

$$
\overline{a} \cdot_n (\overline{b} +_n \overline{c}) = \overline{a} \cdot_n \overline{b} + c
$$

$$
= \overline{a(b+c)}
$$

$$
= \overline{ab} + \overline{ac}
$$

$$
= \overline{ab} +_n \overline{ac}
$$

$$
= \overline{a} \cdot_n \overline{b} +_n \overline{a} \cdot_n \overline{c}.
$$

The commutative and associative laws follow similarly from the commutative laws and associative laws for ^Z.