

Note: in parts c and d I'll accept any plausible argument based on the properties of 'people' and the logical structure of \mathcal{L}_{max} is as a understanding as \mathcal{L}_{max} of the logical structure of the logical str the statement!

(a) (A ∨ B)∧ ∼ (A ∧ B) or (A∧ ∼ B) ∨ (∼ A ∧ B) or (A ∨ B) ∧ (∼ A∨ ∼ B). Other forms are possible, check the truth table. (Fewer marks for unnecessarily complex answers.)

(Table not required in the solution, but recommended.)
(b) $(A \wedge B \wedge \sim B \wedge \sim C) \vee (\sim A \wedge B \wedge \sim C) \vee (\sim A \wedge B \wedge C) \vee (A \wedge B \wedge C)$. A bonus mark for seeing $\begin{array}{c} \hbox{how to simplify this John McK can't-Arg of the truth table.} \end{array}$ how to simplify this John McK can't. Again, check the truth table.

(Table not required in the solution, but recommended.)
(c) $\forall x \in P, \exists y \in P, \exists t \in T \text{ such that } R(x, y, t).$

 $(\mathbf{m}, \mathbf{m}, \mathbf{$ $(1 + \epsilon)$. The common and 'such that' are recommended but not necessary. All the symbols are necessary sary.)

- (d) $\sim (\exists x \in P, \forall y \in P, \forall t \in T \text{ such that } R(x, y, t))$ or (better) $\forall x \in P, \exists y \in P, \exists t \in T \text{ such that } \sim R(x, y, t)$. (Comments as for part c.)
- (e) $\forall a \in \mathbb{Z}, \forall b \in \mathbb{Z}, R(a, b) \Longrightarrow (\exists c \in \mathbb{Q} \text{ such that } R(a, c) \wedge R(c, b)).$

(Comments as for previous 2 parts, but marks lost for not using the symbols \Longrightarrow , $R(.)$ and \land .) Proof: Let $a, b \in \mathbb{Z}$ satisfy $R(a, b)$, in other words $a < b$. Define $c = (a+b)/2$. Clearly $c \in \mathbb{Q}$. We claim that $R(a, c) \wedge R(c, b)$ is true; in other words $a < c$ and $c < b$. $c = (a+b)/2 < (b+b)/2 = b$ since $a < b$, and $a = (a + a)/2 < (a + b)/2 = c$ since $a < b$. So the claim is established, since c satisfies the requirements of the theorem. satisfies the requirements of the theorem.

2. (A ⇒⇒ B)∧ → C ⇒ C ⇒ C)

$$
(A \implies B) \land \sim (A \iff C) \iff (\sim A \lor B) \land \sim ((\sim A \lor C) \land (A \lor \sim C))
$$

$$
\iff (\sim A \lor B) \land (\sim (\sim A \lor C) \lor \sim (A \lor \sim C)) \iff (\sim A \lor B) \land ((A \land \sim C) \lor (\sim A \land C)).
$$

 $\frac{1}{\sqrt{2}}$ approach is to prove the approach is to prove the $\frac{1}{\sqrt{2}}$ and $\frac{1}{\sqrt{2}}$ expanding.
Markers note; only give credit for this if the rules are proven or a source referenced! Not covered

in text or in class. Another way is to read off the expression from the truth table. (One bracket in text or in class. Another way is to read off the expression from the truth table table. (One bracket) for every line that ends in a \mathcal{L} think about it.

$$
(A \land B \land \sim C) \lor (\sim A \land B \land C) \lor (\sim A \land \sim B \land C)
$$

(A ∧ B∧ ∼ C) ∨ (∼ A ∧ B ∧ C) ∨ (∼ A∧ ∼ B ∧ C) $\frac{1}{2}$ $\frac{1}{2}$ using a truth table!)

 $\sim (\sim (A \vee \sim B) \wedge \sim (\sim A \vee B))$ is a tautology, and can be expressed as (for example) $A \vee \sim A$.

3. (a) ⇒

Let $n \in \mathbb{Z}$ be divisible by 3. Then $n = 3k$ for some integer k, and $n^3 = 27k^3 = 3(9k^3)$.
Therefore since $0k^3 \in \mathbb{Z}$ n^3 is divisible by 3 as required Therefore, since $9k^3 \in \mathbb{Z}$, n^3 is divisible by 3 as required.
 \Leftarrow

ं
π⊥∹ It is easiest to prove the contrapositive of the statement, which is σ does not divide n^{3n} does not divide n^{3} "
There are 2 cases to check; $n = 3k + 1$ and $n = 3k + 2$.

 T_{real} (2, + 1)3 $27L^3 + 27L^2 + 0L + 1$ $2(0L^3 + 0L)$ Case1: $(3k + 1)^3 = 27k^3 + 27k^2 + 9k^2 + 1 = 3(9k^3 + 9k^2 + 3k) + 1$ which is not divisible by 3
since $9k^3 + 9k^2 + 3k \in \mathbb{Z}$

since $9k^3 + 9k^2 + 3k \in \mathbb{Z}$.
C₂₀₀ $(9k+9)^3 - 97k^3$ Case2: $(3k + 2)^3 = 21k^3 + 34k^2 + 30k + 8 = 3(9k^3 + 18k^2 + 12k + 2) + 2$ which is not divisible
by 3 since $9k^3 + 18k^2 + 12k + 2 \in \mathbb{Z}$ by 3 since $9k^3 + 18k^2 + 12k + 2 \in \mathbb{Z}$.
(Note that we use the "in exactly one way" part of the facts I said you could use here.)

 \mathcal{L} is the that we use the "including one way" part of the facts I said you could use here.) (b) Assume, for a contradiction, that there is a rational number d that satisfies $a^2 = 3$. Then we may assume that $d - n/a$ where $n, a \in \mathbb{Z}$ have no common factor may assume that $d = p/q$ where $p, q \in \mathbb{Z}$ have no common factor.

Then
$$
3 = d^3 = p^3/q^3
$$
, so

$$
p^3 = 3q^3. (*)
$$

This implies (since $q^* \in \mathbb{Z}$) that p³ is divisible by 3. From part a of this question, we conclude that n is divisible by 3. Let $n = 3t$, where $t \in \mathbb{Z}$

that p is divisible by 3. Let $p = 3t$, where $t \in \mathbb{Z}$.
Substituting in (*) we obtain $3q^3 = 27t^3$ or $q^3 = 3(3t^3)$. This implies that q^3 is divisible by 3 Substituting in (*) we obtain $3q^3 \equiv 27t^3$ or $q^3 \equiv 3(3t^3)$. This implies that q^3 is divisible by 3
and (using part a again) that q is divisible by 3. Thus BOTH n and q are divisible by 3; in other $\frac{1}{2}$ and $\frac{1}{2}$ is divisible by $\frac{1}{2}$ is divisible by $\frac{1}{2}$ and $\frac{1}{2}$ are divisible by $\frac{1}{2}$ are divisible by $\frac{1}{2}$ are divisible by $\frac{1}{2}$ common factor. This gives a contradiction to our o where $\frac{1}{2}$ have a common factor. This gives a contradiction to our original assumption to our original assumptions, and $\frac{1}{2}$ proves the result.

- 4. (a) Prove $A \setminus (A \setminus B) = A \cap B$. We prove $x \in A \setminus (A \setminus B) \iff x \in A \cap B$.
 \iff Let $x \in A \setminus (A \setminus B)$. Then $x \in A$ and $x \notin A \setminus B$. But $A \setminus B = \{x \in A : x \notin B\}$. Therefore we must have $x \in B$. Since x is in both A and B, $x \in A \cap B$ as required. $\mathbf{A} \times \mathbf{B} = \mathbf{A} \times \mathbf{B} \times \mathbf{C}$ is in both $\mathbf{A} \times \mathbf{C}$. Since $\mathbf{A} \times \mathbf{B}$ and $\mathbf{B} \times \mathbf{C}$ and $\mathbf{A} \times \mathbf{C}$. Since $\mathbf{A} \times \mathbf{C}$ $\frac{1}{2}$
	- $\frac{1}{2}$ (b) Prove $(A \cup B^c)^c = A^c \cap B$. Let U be the universal set, and let $x \in U$. We prove $x \in (A \cup B^c)^c \iff$ $x \in A^* \sqcup D.$
 $\qquad \qquad \cdot \quad \mathbf{I} \quad \mathbf{A} \quad \ldots$ \Rightarrow Let $x \in (A \cup B^c)^c$
 $x \in A^c \cap B$ as required . Then $x \notin A \cup B$, so $x \notin A$ and $x \in B$. Since $x \notin A$, $x \in A$ so $x \in A^* \cap D$ as required. \Leftarrow Let $x \in A^* \cap D$. Then $x \in A^*$ and $x \in D$, so $x \notin A$ and $x \notin D^*$. Therefore $x \notin A \cup D^*$
and $x \in (A \cap B^c)^c$ as required , and $x \in (A \cup B^c)$ as required.
	- (c) We need to prove 2 things; that the empty set ϕ has no proper subset, and that every nonempty set has a proper subset.

Assume (for a contradiction) that A is a proper subset of ϕ . Then (by the definition of subset) $x \in A \Longrightarrow x \in \phi$. Since ϕ has no elements we conclude that A has no elements. In other words, $A = \phi$. But the definition of proper subset requires that $A \neq \phi$ giving the contradiction we $\frac{1}{4}$ = $\frac{1}{4}$ 6 $\frac{1}{4}$ 6 $\frac{1}{4}$ 6 $\frac{1}{4}$ 6 $\frac{1}{4}$ 6 $\frac{1}{6}$ 6 $\frac{1}{6}$ require.
Now assume that we have a set A such that $A \neq \phi$. We show that ϕ is a proper subset of A.

To prove that $\phi \subset A$ we need to show that $r \subset \phi \longrightarrow r \subset A$. This is vacuously true since ϕ . To prove that $\frac{1}{\sqrt{2}}$ and $\frac{1}{\sqrt{2}}$ and $\frac{1}{\sqrt{2}}$ and $\frac{1}{\sqrt{2}}$ is variable $\frac{1}{\sqrt{2}}$. This is variable $\frac{1}{\sqrt{2}}$ has no elements. Lastly, ϕ is a proper subset because there must be some x such that $x \in A$ (otherwise A would be the empty set).

 $(N_{\text{data}} + \text{height})$. $\frac{1}{\sqrt{N}}$ is very simple, but needs to be carefully worded to be a water to be a water to be a water $\frac{1}{\sqrt{N}}$ proof.)